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EXAMINER

WONG, LESLIE

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Please find below and/or attached an Office communication concerning this application or proceeding.

The time period for reply, if any, is set in the attached communication.

Office Action Summary	Application No. 10/056,880	Applicant(s) MILBY, GREGORY H.	
	Examiner Leslie Wong	Art Unit 2164	

-- The MAILING DATE of this communication appears on the cover sheet with the correspondence address --

Period for Reply

A SHORTENED STATUTORY PERIOD FOR REPLY IS SET TO EXPIRE 3 MONTH(S) OR THIRTY (30) DAYS, WHICHEVER IS LONGER, FROM THE MAILING DATE OF THIS COMMUNICATION.

- Extensions of time may be available under the provisions of 37 CFR 1.136(a). In no event, however, may a reply be timely filed after SIX (6) MONTHS from the mailing date of this communication.
- If NO period for reply is specified above, the maximum statutory period will apply and will expire SIX (6) MONTHS from the mailing date of this communication.
- Failure to reply within the set or extended period for reply will, by statute, cause the application to become ABANDONED (35 U.S.C. § 133). Any reply received by the Office later than three months after the mailing date of this communication, even if timely filed, may reduce any earned patent term adjustment. See 37 CFR 1.704(b).

Status

- 1) ☒ Responsive to communication(s) filed on 18 April 2005.
- 2a) ☒ This action is **FINAL**. 2b) ☐ This action is non-final.
- 3) ☐ Since this application is in condition for allowance except for formal matters, prosecution as to the merits is closed in accordance with the practice under *Ex parte Quayle*, 1935 C.D. 11, 453 O.G. 213.

Disposition of Claims

- 4) ☒ Claim(s) 1-12, 14, 16-32 is/are pending in the application.
- 4a) Of the above claim(s) _____ is/are withdrawn from consideration.
- 5) ☐ Claim(s) _____ is/are allowed.
- 6) ☒ Claim(s) 1-12, 14, 16-25, 27, 28, 30 and 32 is/are rejected.
- 7) ☒ Claim(s) 26, 29, and 31 is/are objected to.
- 8) ☐ Claim(s) _____ are subject to restriction and/or election requirement.

Application Papers

- 9) ☐ The specification is objected to by the Examiner.
- 10) ☐ The drawing(s) filed on _____ is/are: a) ☐ accepted or b) ☐ objected to by the Examiner.
Applicant may not request that any objection to the drawing(s) be held in abeyance. See 37 CFR 1.85(a).
Replacement drawing sheet(s) including the correction is required if the drawing(s) is objected to. See 37 CFR 1.121(d).
- 11) ☐ The oath or declaration is objected to by the Examiner. Note the attached Office Action or form PTO-152.

Priority under 35 U.S.C. § 119

- 12) ☐ Acknowledgment is made of a claim for foreign priority under 35 U.S.C. § 119(a)-(d) or (f).
- a) ☐ All b) ☐ Some * c) ☐ None of:
1. ☐ Certified copies of the priority documents have been received.
2. ☐ Certified copies of the priority documents have been received in Application No. _____.
3. ☐ Copies of the certified copies of the priority documents have been received in this National Stage application from the International Bureau (PCT Rule 17.2(a)).
- * See the attached detailed Office action for a list of the certified copies not received.

Attachment(s)

- | | |
|--|---|
| 1) <input checked="" type="checkbox"/> Notice of References Cited (PTO-892) | 4) <input type="checkbox"/> Interview Summary (PTO-413)
Paper No(s)/Mail Date. _____ |
| 2) <input type="checkbox"/> Notice of Draftsperson's Patent Drawing Review (PTO-948) | 5) <input type="checkbox"/> Notice of Informal Patent Application |
| 3) <input type="checkbox"/> Information Disclosure Statement(s) (PTO/SB/08)
Paper No(s)/Mail Date _____ | 6) <input type="checkbox"/> Other: _____ |

DETAILED ACTION

Response to Amendment

1. Receipt of Applicant's Amendment, filed 18 April 2005, is acknowledged.

Claim Rejections - 35 USC § 103

2. The following is a quotation of 35 U.S.C. 103(a) which forms the basis for all obviousness rejections set forth in this Office action:

(a) A patent may not be obtained though the invention is not identically disclosed or described as set forth in section 102 of this title, if the differences between the subject matter sought to be patented and the prior art are such that the subject matter as a whole would have been obvious at the time the invention was made to a person having ordinary skill in the art to which said subject matter pertains. Patentability shall not be negated by the manner in which the invention was made.

3. Claims 1-10, 14, 18-25, 27-28, 30, and 32 are rejected under 35 U.S.C. 103(a) as being unpatentable over **Krishna** (U.S. Patent 6,138,111A) in view of **Dessloch** (U.S. Patent 6,338,056B1).

Regarding claim 1, **Krishna** teaches a method of performing a join in a database system comprising:

- a). **'receiving a join query containing a selection predicate and a projection'** as a query containing one or more predicates which specify the information that the query should retrieve from the data and select the names (i.e., projection) of all people who live and work in the same city (col. 1, lines 43-57);
- b). **'determining a cost associated with applying the function on a first table and a cost associated with applying the function on a second table'** as a new

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metric, Sigma, is used to calculate the estimated cost of the joined tables R, S, and T by picking a join order having a smallest Sigma among all join orders (col. 3, lines 44-52); and

c). **'selecting a join path based on relative costs of applying the function on the first and second tables'** as Sigma for join order (1) is 80 and Sigma for join order (2) is 560. Consequently, order (1) is preferred over order (2) because joiner order (1) offers an optimal order for the join tables which consume less time and resources to process the join query (col. 3, lines 52-55 and lines 23-29).

a). **Krishna** does not explicitly teach a step of a join query containing at least one function selected from the group consisting of a selection predicate applied on a *complex attribute*, a projection applied on a complex attribute, and a *user-defined data type* method.

Dessloch, however, teaches **'a join query containing at least one function selected from the group consisting of a selection predicate applied on a complex attribute, a projection applied on a complex attribute, and a user-defined data type method'** as object-relational database which allows relational database systems to store text, images, video, audio, and other non-traditional data types. These databases are extensible in terms of their type system and their query language, thereby allowing the user to create and query new data types of the mentioned contents in the database (col. 2, lines 18-30; col. 5, line 61 – col. 6, line 37). As specified in Applicant's

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Specification at page 5, paragraph 0029, examples of complex data types include data types defined for storing image data, audio data, video data etc.

It would have been obvious to one of ordinary skill in the data processing art at the time of the invention was made to combine the teachings of the cited references because **Dessloch's** teaching would have allowed **Krishna's** to store and search for the new types of contents, such as images, video, and audio, based on non-traditional data types by enabling users to add columns of the appropriate data types as suggested by **Dessloch** at col. 6, lines 1-20. Further, complex attribute and user-defined data types as suggested by **Dessloch** improve the performance of the database system because the optimizer will attempt to take advantage of the User Defined Functions (UDFs data filter) and the index exploitation of the UDFs on a join operation for the generation of the optimal plan (col. 17, lines 40-67).

Regarding claim 2, **Krishna** further teaches wherein **'selecting the join path comprises applying the function on one of the first and second tables associated with a lower cost'** as using the Sigma metric, the join order having the smallest Sigma among all join orders is selected as optimal and used to perform the join. Sigma for join order (1) is 80 and Sigma for join order (2) is 560. Order (1) is preferred over order (2) because order (1) offers an optimal order for the join tables. Thus, the cost for processing the joined query using order (1) would be lower than order (2) (col. 3, lines 47-55 and lines 23-29).

Regarding claim 3, **Krishna** further teaches wherein **'determining the costs comprises determining the respective cardinalities of the first and second tables'** as estimating the cost for a join by broken the join orders and obtain the cardinality of each component join in the join order (col. 3, lines 52-64).

Regarding claim 4, **Krishna** further teaches wherein **'determining the cost of applying the function on the second table comprises determining the cost of a join table that is a result of a join of the first table and another table'** as joining tables R and S, and then join the result with table T, resulted in join order 1 which contains the smallest number of records for the total join. The ties among competing join orders may be broken using some other heuristic to perform cost estimate calculations for the selected join order (col. 3, lines 30-65). Joiner order (1) offers an optimal order for the join tables which consume less time and resources to process the join query (col. 3, lines 52-55 and lines 23-29).

Regarding claim 5, **Krishna** further teaches wherein **'selecting the join path comprises applying the function on one of the first and second tables that has the lower cardinality'** as a value for Sigma is calculated by summing each of the cardinality estimates for each component join in the join order and select the join order with the smallest value of Sigma to perform the query (col. 3, line 66 – col. 4, line 7).

Regarding claim 6, **Krishna** further teaches wherein the function comprises a selection predicate applied on a complex attribute of the first table, the join query further containing a projection applied on a complex attribute of the first table, the method further comprising:

a). **'determining a cost associated with applying the projection on the first table and a cost associated with applying the projection on the join table'** as $R \text{ joins } S = 20$ and $((R \text{ join } S) \text{ join } T) = 60$, then the total number of tuples created during the calculation of the total query is 80. The join order with the smallest Sigma value is selected to perform the join query because it consumes less time and resources to process the query (col. 3, lines 23-42),

b). **'wherein selecting the join path comprises applying the projection on one of the first table and the join table associated with a lower cost'** as join order 1 which contains the smallest number of records for the total join. Thus, offers an optimal join for the tables (col. 3, lines 30-64).

Regarding claim 7, **Krishna** further teaches wherein **'selecting the join path comprises applying the projection on one of the first table and join table with the lower cardinality'** as joining tables R and S, and then join the result with table T, resulted in join order 1 which contains the smallest number of records for the total join (col. 3, lines 30-64).

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Regarding claim 8, **Krishna** further teaches wherein '**identifying the function as a costly function**' as calculating the estimated cost for the selected join order (col. 3, lines 44-52).

Regarding claim 9, **Krishna** further teaches wherein '**the receiving, determining, and selecting acts are performed by an optimizer module**' as a database optimizer may use whatever other method it has available for computing or estimating the selectivity of other kinds of joins (col. 5, lines 32-38).

Regarding claim 10, **Krishna** further teaches wherein '**determining the costs of applying the function on the first and second tables comprises determining the costs of applying the function on relational tables**' as calculating the cardinalities of the selected join order for the tables specified in the join query (col. 3, lines 30-64).

Krishna does not explicitly teach determining the costs of applying the function on *object* relational tables.

Dessloch, however, teaches '**object relational tables**' as object-relational database systems provide an architecture and application program interface for integrating content management and search functions for new data types in form of plug-ins (col. 2, lines 18-39).

It would have been obvious to one of ordinary skill in the data processing art at the time of the invention was made to combine the teachings of the cited references because **Dessloch's** teaching would have allowed **Krishna's** to provide database

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users with new data types and their advanced content search capabilities inside of SQL as suggested by **Dessloch** at col. 2 lines 34-39.

Regarding claim 14, **Krishna** further teaches wherein **'the join query specifies the function being applied on a first table'** as a query containing one or more predicates which specify the information that the query should retrieve from the data and select the names (i.e., projection) of all people who live and work in the same city (col. 1, lines 43-57), and

wherein **'the instructions when executed cause the database system to determine the join path by applying the function on a second table having a lower cardinality than the first table'** as using the Sigma metric, the join order having the smallest Sigma among all join orders is selected as optimal and used to perform the join. Sigma for join order (1) is 80 and Sigma for join order (2) is 560. Order (1) is preferred over order (2). Estimating the cost for a join by broken the join orders and obtain the cardinality of each component join in the join order and select the join order which has the smallest value of cardinalities to perform the join query (col. 3, lines 52-64) (col. 3, lines 47-55).

Regarding claim 18, **Krishna** further teaches a database system comprising:

a). **'a storage system to store tables'** as a mass storage device (Fig. 1, element 16); and

b). **'an optimizer to receiving a join query containing a selection predicate and a projection'** as implementing the methods in query optimizers in relational database management systems. A query containing one or more predicates which specify the information that the query should retrieve from the data and select the names (i.e., projection) of all people who live and work in the same city (col. 1, lines 43-57 and abstract);

c). **'the optimizer to select a join plan based at least in part on a comparison of a first cost of applying the function on a first table and a second cost of applying the function on a second table'** as Sigma for join order (1) is 80 and Sigma for join order (2) is 560. Consequently, order (1) is preferred over order (2) because joiner order (1) offers an optimal order for the join tables which consume less time and resources to process the join query (col. 3, lines 52-55 and lines 23-29).

Krishna does not explicitly teach a step of a join query containing at least one function selected from the group consisting of a selection predicate applied on a *complex attribute*, a projection applied on a complex attribute, and a *user-defined data type* method.

Dessloch, however, teaches **'a join query containing at least one function selected from the group consisting of a selection predicate applied on a complex attribute, a projection applied on a complex attribute, and a user-defined data type method'** as object-relational database which allows relational database systems to

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store text, images, video, audio, and other non-traditional data types. These database are extensible in terms of their type system and their query language, thereby allowing the user to create new data types of contents in the database as well as manipulate and search such content (col. 2, lines 18-30; col. 5, line 61 – col. 6, line 37). As specified in Applicant's Specification at page 5, paragraph 0029, examples of complex data types include data types defined for storing image data, audio data, video data etc.

It would have been obvious to one of ordinary skill in the data processing art at the time of the invention was made to combine the teachings of the cited references because **Dessloch's** teaching would have allowed **Krishna's** to store and search for the new types of contents, such as images, video, and audio, based on non-traditional data types by enabling users to add columns of the appropriate data types as suggested by **Dessloch** at col. 6, lines 1-20. Further, complex attribute and user-defined data types as suggested by **Dessloch** improve the performance of the database system because the optimizer will attempt to take advantage of the User-Defined functions (UDFs data filter) and the index exploitation of the UDFs on a join operation for the generation of the optimal plan (col. 17, lines 40-67).

Regarding claim 19, **Krishna** further teaches wherein **'the optimizer is to select the join plan that applies the function on the one of the first table and second table with a lower cardinality'** as a value for Sigma is calculated by summing each of the cardinality estimates for each component join in the join order and select the join

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order with the smallest value of Sigma to perform the query (col. 3, line 66 – col. 4, line 7).

Regarding claim 20, **Krishna** further teaches wherein **'the second table is a join result of the first table and another table'** as a join query in which tables R, S, and T are to be joined. The system joins tables R and S then joins the result of tables R and S with table T (col. 3, lines 30-41).

Regarding claim 21, **Krishna** further teaches wherein **'the join query specifies the function being applied on the first table and the optimizer to apply the function the second table rather than the first table in response to determining the second cost is lower than the first cost'** as selecting all columns from table S (i.e. first table) (col. 4, lines 40-41; col. 3, line 66 – col. 4, line 7).

Regarding claim 22, **Krishna** further teaches wherein **'the first and second tables are relational tables'** as a general computer platform suitable for supporting a relational database system. A relational database is a database that is perceived by its users as a collection of tables such as tables R, S, and T in Fig. 1 (col. 2, lines 66-67; col. 1, lines 7-9).

Krishna does not explicitly teach **"object relational tables"**.

Dessloch, however, teaches **'the first and second tables are object relational tables'** as object-relational database systems provide an architecture and application program interface for integrating content management and search functions for new data types in form of plug-ins (col. 2, lines 18-39).

It would have been obvious to one of ordinary skill in the data processing art at the time of the invention was made to combine the teachings of the cited references because **Dessloch's** teaching would have allowed **Krishna's** to provide database users with new data types and their advanced content search capabilities inside of SQL as suggested by **Dessloch** at col. 2 lines 34-39.

Regarding claim 23, **Krishna** further teaches wherein **'the join query further specifies application of a second function selected from the group consisting a selection predicate and a projection'** as a query containing one or more predicates which specify the information that the query should retrieve from the data and select the names (i.e., projection) of all people who live and work in the same city (col. 1, lines 43-57), **'the second function being applied on a third table, wherein the instructions when executed cause the database system to determine the join path by further applying the second function on one of the third table and a fourth table with a lower cardinality, wherein 'the fourth table is a join result of the third table and another table'** as a join query in which tables R, S, and T are to be joined. Join tables

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R and S which yields the temp table containing the results of R and S, and then join the result of tables R and S with table T (col. 3, lines 30-41).

Krishna does not explicitly teach a step of a selection predicate applied on a complex attribute, a projection applied on a complex attribute, and a user-defined data type method.

Dessloch, however, teaches ‘a selection predicate applied on a complex attribute, a projection applied on a complex attribute, and a user-defined data type method’ as object-relational database which allows relational database systems to store text, images, video, audio, and other non-traditional data types. These database are extensible in terms of their type system and their query language, thereby allowing the user to create new data types of contents in the database as well as manipulate and search such content (col. 2, lines 18-30; col. 5, line 61 – col. 6, line 37).

Regarding claim 24, **Krishna** further teaches wherein ‘**the tables are relational tables**’ as a general computer platform suitable for supporting a relational database system. A relational database is a database that is perceived by its users as a collection of tables such as tables R, S, and T in Fig. 1 (col. 2, lines 66-67; col. 1, lines 7-9).

Krishna does not explicitly teach “**object relational tables**”.

Dessloch, however, teaches '**the tables are *object relational tables***' as object-relational database systems provide an architecture and application program interface for integrating content management and search functions for new data types in form of plug-ins (col. 2, lines 18-39).

It would have been obvious to one of ordinary skill in the data processing art at the time of the invention was made to combine the teachings of the cited references because **Dessloch's** teaching would have allowed **Krishna's** to provide database users with new data types and their advanced content search capabilities inside of SQL as suggested by **Dessloch** at col. 2 lines 34-39.

Regarding claims 25 and 30, **Krishna** further teaches wherein the second table contains a result of a join between the first table and another table,

Wherein selecting the join path comprises selecting a join path in which at least one of selection and projection is applied on the second table rather than the first table, the method further comprising:

a). receiving a second query specifying a join of the first table and another table, the second query specifying at least one of a selection predicate applied on a non-complex attribute and a projection applied on a non-complex attribute (col. 3, lines 23-32); and

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b). selecting another join path for the second query in which the selection or projection is applied on the first table before performing a join of the first table with another table (col. 3, lines 35-43).

Regarding claims 27, wherein the query specifies application of the function of the first table,

Wherein selecting the join path comprises selecting the join path in which the function is applied on the second table (col. 1, lines 43-57), the second table containing a join result of a join of the first table and another table (col. 3, lines 30-41).

Regarding claims 32, **Krishna** teaches a method of performing a join a a database system, comprising:

a). receiving a join query specifying a join of a first table and a second table and containing at least one of a selection predicate (i.e., R.Pk) and a projection (i.e., R, S, T) (col. 5, lines 9-11, Example 2);

d). wherein a second join path is selected in which the at least one of the selection predicate and projection is applied on the first table before the join in response to determining that the at least one of the selection predicate and the projection is applied on a non-complex attribute (col. 3, lines 47-55 and lines 23-32).

Krishna does not explicitly teach the steps of:

b). selecting join path for the join query in response to determining whether the at least one of the selection predicate and projection is applied on a complex attribute;

c). wherein a first join path is selected in which the at least one of the selection predicate and projection is applied on a join table in response to determining that the at least one of the selection predicate and projection is applied on a join table in response to determining that the at least one of the selection predicate and projection is applied on a complex attribute, the join table containing a join result of the first and second tables.

Dessloch teaches steps b and c as object-relational database which allows relational database systems to store text, images, video, audio, and other non-traditional data types. These databases are extensible in terms of their type system and their query language, thereby allowing the user to create and query new data types of the mentioned contents in the database (col. 2, lines 18-30; col. 5, line 61 – col. 6, line 37). As specified in Applicant's Specification at page 5, paragraph 0029, examples of complex data types include data types defined for storing image data, audio data, video data etc.

It would have been obvious to one of ordinary skill in the data processing art at the time of the invention was made to combine the teachings of the cited references because **Dessloch's** teaching would have allowed **Krishna's** to store and search for the new types of contents, such as images, video, and audio, based on non-traditional data types by enabling users to add columns of the appropriate data types as

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suggested by **Dessloch** at col. 6, lines 1-20. Further, complex attribute and user-defined data types as suggested by **Dessloch** improve the performance of the database system because the optimizer will attempt to take advantage of the User Defined Functions (UDFs data filter) and the index exploitation of the UDFs on a join operation for the generation of the optimal plan (col. 17, lines 40-67).

4. Claims 11-12, 16-17, and 28 are rejected under 35 U.S.C. 103(a) as being unpatentable over **Krishna** (U.S. Patent 6,138,111A) in view of **Dessloch** (U.S. Patent 6,338,056B1) as applied to claims 1 and further in view of **Ross et al.** ("Ross") (US 6026390 A).

Regarding claim 11, **Krishna** teaches an article comprising at least one storage medium containing instructions that when executed cause a database system to:

a). **'receiving a join query containing a selection predicate, a projection, and a user-defined data type method'** as a query containing one or more predicates which specify the information that the query should retrieve from the data and select the names (i.e., projection) of all people who live and work in the same city (col. 1, lines 43-57); and

b). **'determine a join path for the join query based at least in part on a cost associated with application of the function'** as a new metric, Sigma, is used to calculate the estimated cost by picking a join order having a smallest Sigma among all join orders (col. 3, lines 44-52).

a). **Krishna** does not explicitly teach a step of a join query containing at least one function selected from the group consisting of a selection predicate applied on a *complex attribute*, a projection applied on a complex attribute, and a *user-defined data type* method.

b). **Krishna** does not explicitly teach the *complex attribute*.

Dessloch, however, teaches ‘a join query containing at least one function selected from the group consisting of a selection predicate applied on a complex attribute, a projection applied on a complex attribute, and a user-defined data type method’ as object-relational database which allows relational database systems to store text, images, video, audio, and other non-traditional data types. These database are extensible in terms of their type system and their query language, thereby allowing the user to create new data types of contents in the database as well as manipulate and search such content (col. 2, lines 18-30; col. 5, line 61 – col. 6, line 37). As specified in Applicant’s Specification at page 5, paragraph 0029, examples of complex data types include data types defined for storing image data, audio data, video data etc.

It would have been obvious to one of ordinary skill in the data processing art at the time of the invention was made to combine the teachings of the cited references because **Dessloch’s** teaching would have allowed **Krishna’s** to store and search for

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the new types of contents, such as images, video, and audio, based on non-traditional data types by enabling users to add columns of the appropriate data types as suggested by **Dessloch** at col. 6, lines 1-20. Further, complex attribute and user-defined data types as suggested by **Dessloch** improve the performance of the database system because the optimizer will attempt to take advantage of the User-Defined Functions (UDFs data filter) and the index exploitation of the UDFs on a join operation for the generation of the optimal plan (col. 17, lines 40-67).

Krishna and Dessloch do not explicitly teach the steps of:

- a). the join query specifying a join of a first table and a second table to produce a join table;
- c). wherein determining the join path comprises selecting the join path in which the function is applied on the join table rather than the first table or second table to reduce cost.

Ross, however, teaches step a and "selecting the join path in which the function is applied on the join table" (col. 12, lines 30-55).

It would have been obvious to one of ordinary skill in the data processing art at the time of the invention was made to combine the teachings of the cited references because **Ross'** teaching would have allowed **Krishna- Dessloch's** it can be used to reduce the cost of query optimization by mining the cost for each update relation (col. 12, lines 30-32).

Regarding claim 12, **Krishna** further teaches wherein **‘the join query specifies the function being applied on the first table, and wherein the instructions when executed cause the database system to determine the join path in which the function is applied on the table’** as optimizing the order in which tables are joined by selecting the join path which contains the smallest Sigma value to perform the query in a multiple join query (col. 2, lines 5-12; col. 3, lines 44-56).

Krishna and **Dessloch** do not explicitly teach to determine the join path in which the function is applied on the join table.

Ross, however, teaches to determine the join path in which the function is applied on the join table (col. 12, lines 30-55).

Regarding claim 16, **Krishna** further teaches wherein **‘the instructions when executed cause the system to determine the join path by further specifying a join of the join table and a third table to produce a fourth table’** and then join the result of tables R and S with table T (col. 3, lines 30-41).

Regarding claim 17, **Krishna** further teaches wherein **‘the join query further specifies application of a second function selected from the group consisting a selection predicate and a projection’** as a query containing one or more predicates which specify the information that the query should retrieve from the data and select the names (i.e., projection) of all people who live and work in the same city (col. 1, lines 43-57), **‘the second function being applied on a third table, wherein the instructions**

when executed cause the database system to determine the join path by further applying the second function on one of the third table and a fourth table with a lower cardinality' as a join query in which tables R, S, and T are to be joined. The system join tables R and S which yields the results of R and S then join the result of tables R and S with table T (col. 3, lines 30-41). The value for Sigma is calculated by summing each of the cardinality estimates for each component join in the join order and select the join order with the smallest value of Sigma to perform the query (col. 3, line 66 – col. 4, line 7).

Krishna does not explicitly teach a step of at least one function selected from the group consisting of a selection predicate applied on a *complex attribute*, a projection applied on a complex attribute, and a *user-defined data type* method.

Dessloch, however, teaches '**a selection predicate applied on a complex attribute, a projection applied on a complex attribute, and a user-defined data type method'** as object-relational database which allows relational database systems to store text, images, video, audio, and other non-traditional data types. These database are extensible in terms of their type system and their query language, thereby allowing the user to create new data types of contents in the database as well as manipulate and search such content (col. 2, lines 18-30; col. 5, line 61 – col. 6, line 37).

Regarding claim 28, **Krishna** further teaches wherein the second table contains a result of a join between the first table and another table,

Wherein selecting the join path comprises selecting a join path in which at least one of selection and projection is applied on the second table rather than the first table, the method further comprising:

a). receiving a second query specifying a join of the first table and another table, the second query specifying at least one of a selection predicate applied on a non-complex attribute and a projection applied on a non-complex attribute (col. 3, lines 23-32); and

b). selecting another join path for the second query in which the selection or projection is applied on the first table before performing a join of the first table with another table (col. 3, lines 35-43).

Allowable Subject Matter

5. Claims 26, 29, and 31 are objected to as being dependent upon a rejected base claim, but would be allowable if rewritten in independent form including all of the limitations of the base claim and any intervening claims.

The following is a statement of reasons for the indication of allowable subject matter: Prior art of record fails to teach a combination of elements including an optimizer module performing N-lookahead join planning in which costs for different

combinations of joins of N+2 tables are determined, where N is greater than or equal to one as recited in dependent claim 26, 29, and 31.

Response to Argument

6. Applicant's arguments filed 18 April 2005 have been fully considered but they are not persuasive.

Applicant argues that although Krishma may teach join order selection for a query that specifies functions applied on simple predicates, there is no suggestion anywhere within Krishma of the determining and selecting acts associated with applying the recited on first and second tables as recited in claim 1.

In response to the preceding arguments, Examiner respectfully submits that there is no requirement that an "express, written motivation to combine must appear in prior art references before a finding of obviousness." See *Ruiz v. A.B. Chance Co.*, 357 F.3d 1270, 1276, 69 USPQ2d 1686, 1690 (Fed. Cir. 2004). For example, motivation to combine prior art references may exist in the nature of the problem to be solved (*Ruiz* at 1276, 69 USPQ2d at 1690) or the knowledge of one of ordinary skill in the art (*National Steel Car v. Canadian Pacific Railway Ltd.*, 357 F.3d 1319, 1338, 69 USPQ2d 1641, 1656 (Fed. Cir. 2004)). See MPEP § 2143.01 for a discussion of proper motivation to combine references.

Applicant's amendment necessitated the new ground(s) of rejection presented in this Office action. Accordingly, **THIS ACTION IS MADE FINAL**. See MPEP § 706.07(a). Applicant is reminded of the extension of time policy as set forth in 37 CFR 1.136(a).

A shortened statutory period for reply to this final action is set to expire **THREE MONTHS** from the mailing date of this action. In the event a first reply is filed within **TWO MONTHS** of the mailing date of this final action and the advisory action is not mailed until after the end of the **THREE-MONTH** shortened statutory period, then the shortened statutory period will expire on the date the advisory action is mailed, and any extension fee pursuant to 37 CFR 1.136(a) will be calculated from the mailing date of the advisory action. In no event, however, will the statutory period for reply expire later than **SIX MONTHS** from the date of this final action.

Any inquiry concerning this communication or earlier communications from the examiner should be directed to Leslie Wong whose telephone number is (571) 272-4120. The examiner can normally be reached on Monday to Friday 9:30am - 6:30 pm.

If attempts to reach the examiner by telephone are unsuccessful, the examiner's supervisor, CHARLES RONES can be reached on (571) 272-4085. The fax phone number for the organization where this application or proceeding is assigned is 571-273-8300.

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Leslie Wong
Primary Patent Examiner
Art Unit 2164

LW
September 3, 2007